#### 3 The Mathematics of Sharing

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#### **Lone-Divider Method**

The first important breakthrough in the mathematics of fair division came in 1943, when Steinhaus came up with a clever way to extend some of the ideas in the dividerchooser method to the case of three players, one of whom plays the role of the divider and the other two who play the role of choosers. Steinhaus' approach was subsequently generalized to any number of players N (one divider and N-1 choosers) by Princeton mathematician Harold Kuhn.

#### **Lone-Divider Method for Three Players**

#### **Preliminaries**

One of the three players will be the divider; the other two players will be choosers. Since it is better to be a chooser than a divider, the decision of who is what is made by a random draw (rolling dice, drawing cards from a deck, etc.). We'll call the divider D and the choosers  $C_1$  and  $C_2$ .

#### **Lone-Divider Method for Three Players**

#### Step 1 (Division)

The divider D divides the cake into three shares ( $s_1$ ,  $s_2$ , and  $s_3$ ). D will get one of these shares, but at this point does not know which one. (Not knowing which share will be his is critical – it forces D to divide the cake into three shares of equal value.)

### Lone-Divider Method for Three Players Step 2 (Bidding)

C₁ declares (usually by writing on a slip of paper) which of the three pieces are fair shares to her. Independently,  $C_2$  does the same. These are the bids. A chooser's bid must list every single piece that he or she considers to be a fair share (i.e., worth onethird or more of the cake)-it may be tempting to bid only for the very best piece, but this is a strategy that can easily backfire. To preserve the privacy requirement, it is important that the bids be made independently, without the choosers being privy to each other's bids.

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# **Lone-Divider Method for Three Players Step 3 (Distribution)**

Who gets which piece? The answer, of course, depends on which pieces are listed in the bids. For convenience, we will separate the pieces into two types: C-pieces (these are pieces *chosen* by either one or both of the choosers) and *U*-pieces (these are unwanted pieces that did not appear in either of the bids).

#### **Lone-Divider Method for Three Players**

#### Step 3 (Distribution) continued

Expressed in terms of value, a *U*-piece is a piece that *both* choosers value at less than 33 1/3% of the cake, and a *C*-piece is a piece that *at least* one of the choosers (maybe both) value at 33 1/3% or more. Depending on the number of *C*-pieces, there are two separate cases to consider.

### **Lone-Divider Method for Three Players Case 1**

When there are two or more C-pieces, there is always a way to give each chooser a different piece from among the pieces listed in her bid. (The details will be covered in Examples 3.2 and 3.3.) Once each chooser gets her piece, the divider gets the last remaining piece. At this point every player has received a fair share, and a fair division has been accomplished.

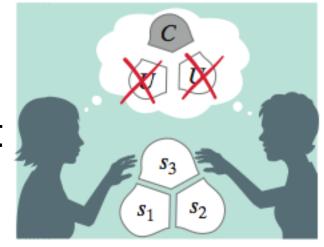
### **Lone-Divider Method for Three Players Case 1**

(Sometimes we might end up in a situation in which  $C_1$  likes  $C_2$ 's piece better than her own and vice versa. In that case it is perfectly reasonable to add a final, informal step and let them swap pieces—this would make each of them happier than they already were, and who could be against that?)

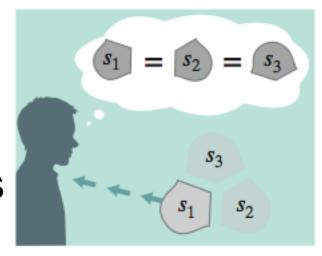
### Lone-Divider Method for Three Players

#### Case 2

When there is only one *C*-piece, we have a bit of a problem because it means that both choosers are bidding for the very same piece.



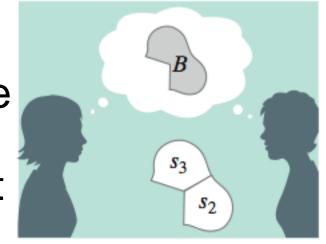
The solution requires a little more creativity. First, we take care of the divider *D*—to whom all pieces are equal in value—by giving him one of the pieces that neither chooser wants.



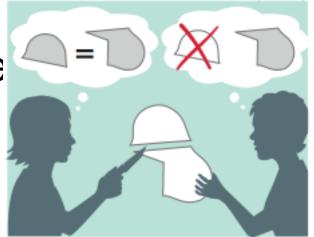
### **Lone-Divider Method for Three Players**

#### Case 2

After *D* gets his piece, the two pieces left (the *C*-piece and the remaining *U*-piece) are recombined into one piece that we call the *B*-piece.



We can revert to the dividerchooser method to finish the fair division: one player cuts the *B*-piece into two pieces; the other player chooses the piece she likes better.



### **Lone-Divider Method for Three Players Case 2**

This process results in a fair division of the cake because it guarantees fair shares for all players. We know that D ends up with a fair share by the very fact that D did the original division, but what about  $C_1$  and  $C_2$ ? The key observation is that in the eyes of both  $C_1$  and  $C_2$  the B-piece is worth more than two-thirds of the value of the original cake (think of the B-piece as 100% of the original cake minus a *U*-piece worth less than 33 1/3%), so when we divide it fairly into two shares, each party is guaranteed more than one-third of the original cake.

Dale, Cindy, and Cher are dividing a cake using Steinhaus's lone-divider method. They draw cards from a well-shuffled deck of cards, and Dale draws the low card (bad luck!) and has to be the divider.

#### Step 1 (Division)

Dale divides the cake into three pieces  $s_1$ ,  $s_2$ , and  $s_3$ . Table 3-1 shows the values of the three pieces in the eyes of each of the players.

TABLE 3-1				
	s <sub>1</sub>	$s_2$	<b>s</b> <sub>3</sub>	
Dale	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	
Cindy	35%	10%	55%	
Cher	40%	25%	35%	

#### Step 2 (Bidding)

We can assume that Cindy's bid list is  $\{s_1, s_3\}$  and Cher's bid list is also  $\{s_1, s_3\}$ .

#### **Step 3 (Distribution)**

The C-pieces are  $s_1$  and  $s_3$ . There are two possible distributions. One distribution would be: Cindy gets  $s_1$ , Cher gets  $s_3$ , and Dale gets  $s_2$ . An even better distribution (the *optimal* distribution) would be: Cindy gets  $s_3$ , Cher gets  $s_1$ , and Dale gets  $s_2$ . In the case of the first distribution, both Cindy and Cher would benefit by swapping pieces, and there is no rational reason why they would not do so.

Step 3 (Distribution) continued

Thus, using the rationality assumption, we can conclude that in either case the final result will be the same: Cindy gets  $s_3$ , Cher gets  $s_1$ , and Dale gets  $s_2$ .

We'll use the same setup as in Example 3.2—Dale is the divider, Cindy and Cher are the choosers.

#### Step 1 (Division)

Dale divides the cake into three pieces  $s_1$ ,  $s_2$ , and  $s_3$ . Table 3-2 shows the values of the three pieces in the eyes of each of the players.

TABLE 3-2				
	s <sub>1</sub>	$s_2$	<b>s</b> <sub>3</sub>	
Dale	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	
Cindy	30%	40%	30%	
Cher	60%	15%	25%	

#### Step 2 (Bidding)

Here Cindy's bid list is  $\{s_2\}$  only, and Cher's bid list is  $\{s_1\}$  only.

#### **Step 3 (Distribution)**

This is the simplest of all situations, as there is only one possible distribution of the pieces: Cindy gets  $s_2$ , Cher gets  $s_1$ , and Dale gets  $s_3$ .

The gang of Examples 3.2 and 3.3 are back at it again.

#### Step 1 (Division)

Dale divides the cake into three pieces  $s_1$ ,  $s_2$ , and  $s_3$ . Table 3-3 shows the values of the three pieces in the eyes of each of the players.

TABLE 3-3				
	s <sub>1</sub>	s <sub>2</sub>	$s_3$	
Dale	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	$33\frac{1}{3}\%$	
Cindy	20%	30%	50%	
Cher	10%	20%	70%	

#### Step 2 (Bidding)

Here Cindy's and Cher's bid list consists of just  $\{s_3\}$ .

#### **Step 3 (Distribution)**

The only C-piece is  $s_3$ . Cindy and Cher talk it over, and without giving away any other information agree that of the two *U*-pieces, s<sub>1</sub> is the least desirable, so they all agree that Dale gets  $s_1$ . (Dale doesn't care which of the three pieces he gets, so he has no rational objection.) The remaining pieces ( $s_2$  and  $s_3$ ) are then recombined to form the B-piece, to be divided between Cindy and Cher using the divider-chooser method (one of them divides

Step 3 (Distribution) continued the B-piece into two shares, the other one chooses the share she likes better). Regardless of how this plays out, both of them will get a very healthy share of the cake: Cindy will end up with a piece worth at least 40% of the original cake (the B-piece is worth 80% of the original cake to Cindy), and Cher will end up with a piece worth at least 45% of the original cake (the B-piece is worth 90% of the original cake to Cher).

The first two steps of Kuhn's method are a straightforward generalization of Steinhaus's lone-divider method for three players, but the distribution step requires some fairly sophisticated mathematical ideas and is rather difficult to describe in full generality, so we will only give an outline here and will illustrate the details with a couple of examples for N = 4 players.

#### **Preliminaries**

One of the players is chosen to be the divider D, and the remaining N-1 players are all going to be choosers. As always, it's better to be a chooser than a divider, so the decision should be made by a random draw.

#### Step 1 (Division)

The divider D divides the set S into N shares  $s_1, s_2, s_3, ..., s_N$ . D is guaranteed of getting one of these shares, but doesn't know which one.

#### Step 2 (Bidding)

Each of the N-1 choosers independently submits a bid list consisting of every share that he or she considers to be a fair share (i.e., worth 1/Nth or more of the booty S).

#### **Step 3 (Distribution)**

The bid lists are opened. Much as we did with three players, we will have to consider two separate cases, depending on how these bid lists turn out.

#### Case 1.

If there is a way to assign a different share to each of the N-1 choosers, then that should be done. (Needless to say, the share assigned to a chooser should be from his or her bid list.) The divider, to whom all shares are presumed to be of equal value, gets the last unassigned share. At the end, players may choose to swap pieces if they want.

#### Case 2.

There is a *standoff*—in other words, there are two choosers both bidding for just one share, or three choosers bidding for just two shares, or K choosers bidding for less than K shares. This is a much more complicated case, and what follows is a rough sketch of what to do. To resolve a standoff, we first set aside the shares involved in the standoff from the remaining shares.

#### Case 2.

Likewise, the players involved in the standoff are temporarily separated from the rest. Each of the remaining players (including the divider) can be assigned a fair share from among the remaining shares and sent packing. All the shares left are recombined into a new booty S to be divided among the players involved in the standoff, and the process starts all over again.

We have one divider, Demi, and three choosers, Chan, Chloe, and Chris.

#### Step 1 (Division)

Demi divides the cake into four shares  $s_1$ ,  $s_2$ ,  $s_3$ , and  $s_4$ . Table 3-4 shows how each of the players values each of the four shares. Remember that the information on each row of Table 3-4 is private and known only to that player.

TABLE 3-4				
	s <sub>1</sub>	s <sub>2</sub>	$s_3$	s <sub>4</sub>
Demi	25%	25%	25%	25%
Chan	30%	20%	35%	15%
Chloe	20%	20%	40%	20%
Chris	25%	20%	20%	35%

#### Step 2 (Bidding)

Chan's bid list is  $\{s_1, s_3\}$ ; Chloe's bid list is  $\{s_3\}$  only; Chris's bid list is  $\{s_1, s_4\}$ .

#### Step 3 (Distribution)

The bid lists are opened. It is clear that for starters Chloe must get  $s_3$  – there is no other option. This forces the rest of the distribution:  $s_1$  must then go to Chan, and  $s_4$  goes to Chris. Finally, we give the last remaining piece,  $s_2$ , to Demi.

This distribution results in a fair division of the cake, although it is not entirely "envy-free" Chan wishes he had Chloe's piece (35% is better than 30%) but Chloe is not about to trade pieces with him, so he is stuck with  $s_1$ . (From a strictly rational point of view, Chan has no reason to gripe—he did not get the best piece, but got a piece worth 30% of the total, better than the 25% he is entitled to.)

Once again, we will let Demi be the divider and Chan, Chloe, and Chris be the three choosers (same players, different game).

#### Step 1 (Division)

Demi divides the cake into four shares  $s_1$ ,  $s_2$ ,  $s_3$ , and  $s_4$ . Table 3-5 shows how each of the players values each of the four shares.

TABLE 3-5				
	s <sub>1</sub>	s <sub>2</sub>	$s_3$	s <sub>4</sub>
Demi	25%	25%	25%	25%
Chan	20%	20%	20%	40%
Chloe	15%	35%	30%	20%
Chris	22%	23%	20%	35%

#### Step 2 (Bidding)

Chan's bid list is  $\{s_4\}$ ; Chloe's bid list is  $\{s_2, s_3\}$  only; Chris's bid list is  $\{s_4\}$ .

#### Step 3 (Distribution)

The bid lists are opened, and the players can see that there is a standoff brewing on the horizon-Chan and Chris are both bidding for  $s_4$ . The first step is to set aside and assign Chloe and Demi a fair share from  $s_1$ ,  $s_2$ , and  $s_4$ . Chloe could be given either  $s_2$  or  $s_3$ . (She would rather have  $s_2$ , of course, but it's not for her to decide.)

#### Step 3 (Distribution)

A coin toss is used to determine which one. Let's say Chloe ends up with  $s_3$  (bad luck!). Demi could be now given either  $s_1$  or  $s_2$ . Another coin toss, and Demi ends up with  $s_1$ . The final move is ... you guessed it!recombine  $s_2$  and  $s_4$  into a single piece to be divided between Chan and Chris using the divider-chooser method.

Step 3 (Distribution)

Since  $(s_2 + s_4)$  is worth 60% to Chan and 58% to Chris (you can check it out in Table 3-5), regardless of how this final division plays out they are both guaranteed a final share worth more than 25% of the cake.

Mission accomplished! We have produced a fair division of the cake.